



arm

TF-M OTP HAL Proposal

TF-M platform layer doesn't support real-world usecases well

- Large amounts of hardcoded crypto keys / attestation values
- Dummy implementation of a lot of key functionality require reimplementing
 - Some platforms have done this, notably for the NV counters which has an almost-working implementation
- Key derivation function returns hardcoded keys
- We don't currently offer any support / guidance on how to implement this as a platform
 - Therefore, no platforms (at least in upstream code) provide a production-ready implementation
 - Our best idea currently is to load keys / attestation values from ITS
 - But we don't have any way to provision ITS, or to create ITS filesystems that can be flashed to devices

CryptoCell provisioning has some problems

- Platforms with a CryptoCell-312 accelerator do support provisioning
- There are some issues with the current system however
 - They are provisioned with random IAKs, which causes problems with testing
 - It only provisions the IAK, HUK, and BL2 ROTPKs
 - And only supports 2 BL2 images, which causes issues with certain platforms (notably diphda)
 - The provisioning system isn't easily extensible to add new data
 - The platforms are still using dummy implementations of the NV counters etc

Proposed solution: add OTP HAL

- Add HAL to allow the platform layer to retrieve data that has been stored in OTP
- Use OTP to store all data that needs to vary on a per-device basis
- Add a mechanism to provision this data on first boot
 - This would then be common across all TF-M platforms, as it uses the OTP HAL to be generic
 - This would be tied to lifecycle-state, which would then also have a real implementation
- By default, use secure on-chip flash instead of real OTP on platforms that lack it
 - This is acceptable to be used in production, provided the flash is locked to secure-privileged access only. Real OTP is still recommended if available however.
 - This makes the OTP HAL compatible with all platforms that support ITS.
 - We can use a tweak in the driver to make this semantically compatible with real OTP (where it's not possible to change a 1 bit to a 0 bit).

Incidental related upgrades

- Since the OTP HAL allows platforms to store secrets in a production-ready way, there is only a small amount non-production-ready code left in the platform layer. Therefore, it makes sense to also:
 - Upgrade the NV counters implementation to one that can handle asynchronous power loss, which makes it acceptable for production
 - ~~Tweak the NV seed implementation be used to provide entropy by default~~ this change went into 1.4, we just need to load the initial seed value from OTP instead of hardcoding it
 - Change the key-derivation operation to perform actual key-derivation instead of returning a hardcoded key
- This leaves us with a platform layer that is by default able to be used in production
 - Provided the platform has internal flash

Design decisions

- API design
- NV counters
- Internal flash allocation
- Provisioning

API design

- API needs to be generic to support different OTP implementations
- We can't make any assumptions on data layout
 - Some implementations may store certain fields in hardware (e.g. CC312 with LCS)
 - Some might have non-contiguous OTP address spaces
- Proposal: Generic getter / setter API
 - Use IDs to indicate which OTP data element to get/set
- Design decision: Don't define how large OTP data elements are
 - Provide an API to get the size of the data element from a given implementation
 - Pro: More flexible – CC312 has limited space so can use more limited NV counters
 - Con: Makes the API harder to use, particularly with regard to memory allocation
 - Alternative: Define sizes in the API

API design (cont.)

```
enum tfm_otp_element_id_t {
    PLAT_OTP_ID_HUK = 0,
    PLAT_OTP_ID_IAK,
    PLAT_OTP_ID_IAK_LEN,
    PLAT_OTP_ID_IAK_TYPE,
    PLAT_OTP_ID_IAK_ID,

    PLAT_OTP_ID_BOOT_SEED,
    PLAT_OTP_ID_LCS,
    PLAT_OTP_ID_IMPLEMENTATION_ID,
    PLAT_OTP_ID_HW_VERSION,
    PLAT_OTP_ID_VERIFICATION_SERVICE_URL,
    PLAT_OTP_ID_PROFILE_DEFINITION,

    < . . . >

    PLAT_OTP_ID_ENTROPY_SEED,
};
```

```
enum tfm_plat_err_t tfm_plat_otp_init(void);

enum tfm_plat_err_t tfm_plat_otp_read(enum tfm_otp_element_id_t id,
                                      size_t out_len, uint8_t *out);

enum tfm_plat_err_t tfm_plat_otp_write(enum tfm_otp_element_id_t id,
                                       size_t in_len,
                                       const uint8_t *in);

enum tfm_plat_err_t tfm_plat_otp_get_size(enum tfm_otp_element_id_t id,
                                           size_t *size);
```

NV counters

- Ideally, we'd use the OTP api for all NV counters
 - However, due to OTP semantics only unary NV counters can be stored in OTP
 - 0x3 is encoded as 0b111, not as 0b011, since we can't unset bits
 - Because of Unary representation OTP NV counters are limited on size
 - 4 bytes gets a counter that goes up to 32
- OTP NV counters are suitable for BL2 images
 - Since they need to count to ~512
- OTP NV counters are not suitable for PS
 - Since we need to be able to do arbitrary amounts of writes, which would require too much OTP space
- Proposal: Use the current NV counter implementation for PS and use the OTP api for BL2
 - Con: We need two distinct code-paths
 - Pro: BL2 counters are more secure (when the platform has real OTP)
 - Alternative: Use the current NV counter implementation for both

Internal flash allocation

- Many platforms have very limited internal flash
- Internal flash backing for OTP HAL requires at least two sectors
 - (not needed if the platform has real OTP)
 - To allow for a backup sector to counter asynchronous power loss during writes
- Proposal: Share space with the NV counters
 - Pro: These also require a backup sector
 - Pro: Can share implementation for restoring from backup

Provisioning



- Provisioning data is input to the device
 - Currently using a debugger, or combining a binary into the image to flash
 - This needs consultation to align best with what market require
- The provisioning service interprets a predefined datastructure, containing all the required assets
- The provisioning service calls the OTP HAL to insert the data into either real OTP or OTP emulated in embedded flash.

Provisioning (cont.)

- Provisioning is done in two stages, corresponding to the two provisioning lifecycle states
 - Currently HUK provisioning is done in the first stage, all other in the second stage
 - This may be changed based on feedback
 - Provisioning writes to the LCS OTP ID to perform a state change
 - How this is handled may depend on the hardware (Notably CC312)
- Provisioning is done via the OTP HAL
 - This means it is the same on all platforms
- By default, a dummy provisioning bundle will be loaded which contains the current dummy TF-M keys / data
- Where possible, provisioning will continue to run TF-M afterwards
 - Not on CC312, as it requires reboots between lifecycle states
 - This allows a first-run experience that is the same as the current state

Provisioning (cont.)

- Design decision: How to inject the keys/provisioning data
 - Currently there is an empty struct, the provisioner can fill the struct with keys via a debugger
 - The struct contains a magic value that must be set in order for it to be used for provisioning
 - Possibility: A HAL function to fill the struct
 - Allows platforms to have flexibility
 - Possibility: Provide default implementation of a serial-port protocol
 - Read data from the UART?
 - It's unclear how well these fit into actual provisioning flows
 - Feedback would be welcomed
 - Does this flow fit with intermediate environments (IDEs etc)?
 - Does this flow fit with factory provisioning flows?

Dummy keys

- Dummy keys are still used during development
- There is now a runtime warning if dummy keys are provisioned to the device
- There is also a runtime warning if the keys on the device are dummy ones
 - In case the code is reflashed to disable dummy provisioning, but the OTP isn't / can't be reset
 - This is done by comparing with the first 32 bytes of the dummy HUK

Future improvements

- Support for encrypted provisioning data
 - Unclear which keys would be used for provisioning
- Support for temporary provisioning LCSes with secure debug disabled
 - To prevent leaking decrypted keys
- On-device key generation where entropy is available

Any Questions?

- Patches at <https://review.trustedfirmware.org/c/TF-M/trusted-firmware-m/+10595/2>
 - Subject to changes from feedback

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Thank You

Danke

Gracias

谢谢

ありがとう

Asante

Merci

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