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TF-M Library model vs. TF-M SFN model

Based on a quantitative approach

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Purpose

- Compare TF-M Library model and SFN model implementation based on some quantitative data
 - Memory footprint
 - Performance
 - Development/management effort
- + Implementation details/functionalities are not the focus
 - Refer to <u>FF-M 1.1 extension</u> Appendix C for detailed analysis on these existing frameworks

TF-M Library model

+ A TF-M specific lightweight framework implementation

- Secure Processing Environment (SPE) as a *secure library*
 - + Based around a set of secure service functions
 - + Those functions run as callbacks from Secure Partition Manager (SPM)
 - + Each secure service function exports its corresponding veneer function
- Use cases
 - + Isolation level 1
 - + Highly resource-constrained devices
 - + Single Armv8-M TrustZone scenario



Secure Function (SFN) Model

-- A new framework defined in FF-M 1.1 extensions

- A simpler programming model compared to IPC model
 - + Reuse RoT service secure function call concept from TF-M Library model and integrate this into FF-M
 - RoT services are implemented as Secure Functions (SFN) that are called by the framework when the client makes a
 request to the service
 - + Identical PSA Client APIs with IPC model
 - + Reduce framework overhead for systems that do not require high levels of isolation
- Use cases
 - + Isolation level 1
 - + Highly resource-constrained devices



- + Memory footprint
 - Similar flash consumption
 - + Code + RO data + RW data
 - SFN model consumes less RAM than Library model does
 - + RW data + ZI data







* Test config: Profile Small, Armclang 6.18, AN521, MinSizeRel build type



+ SFN model costs longer in client calls than Library model does

- Additional operations required by FF-M
 - + Client permission verification
 - + RoT service version validation
 - + Input parameter overlapping checks to avoid double-fetch inconsistency
 - + Message construction and parse
 - + RoT service invokes psa_read() to read input parameters





*TF-M Profiler tool *Test config: Profile Small GNULArm Mu

-- SFN model costs longer in client calls than Library model does (cont'd)

- TF-M specific implementation of SFN model
 - + SFN models shares some common routines with IPC model to simplify implementation/maintenance
 - Such as dynamic handle/message instance allocation
 - Can be optimized further if required
 - + Entry functions for TF-M stateless RoT services
 - TF-M Secure Partition implement an entry function to dispatch stateless RoT service callbacks
 - Reduce consumption of stateless handles to reserve indexes for 3rd-party RoT service usage
 - 3rd-party RoT services can export RoT service callbacks directly without an entry function



-- Development/Maintenance effort

- How many conditional checks/branches are maintained for Library mode/SFN model?
 - + Each one wraps Library/SFN model specific implementation in shared routines with IPC model
 - Changes of Library/SFN may impact IPC model, and vice versa
 - "Bidirectional" development/maintenance effort with IPC model
 - + Library model: TFM_PSA_API/TFM_LIB_MODEL
 - Dedicated standalone SPM/HAL implementation
 - + SFN model: CONFIG_TFM_PSA_API_SFN_CALL/CONFIG_TFM_SPM_BACKEND_SFN
 - Share common routines/implementation with IPC model

	Library model	SFN model
C code	114	9
Linker scripts	17	0
Build system (including manifest tool)	74	11
Total	215	20

#elif CONFIG_TFM_PSA_API_SFN_CALL == 1
\$<\$<BOOL:\${CONFIG_TFM_SPM_BACKEND_SFN}>:...>

- Observations
 - Similar memory footprint
 - SFN model is "slower" due to more execution steps compliant with FF-M
 - Less development/maintenance effort for SFN model, with essential IPC model

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